# Drumul spre Inteligența artificială de încredere trece (și) prin Teoria Jocurilor ?

(și despre echilibre Kantiene)

Gabriel Istrate Universitatea din București

gabrielistrate@acm.org



# Ce caută un teoretician la un Workshop de Inteligență Artificială, la Diaspora Științifică ?

- Am fost "diaspora" 13 ani (Rochester, Los Alamos).
- Cariera ştiinţifică "de bază": algoritmi şi complexitate (dar şi sisteme complexe, simulări sociale multiagent).
- Recent: teoria algoritmică a jocurilor și sisteme multiagent: lucrări (long/BlueSky) la AAMAS in fiecare an 2019-2022, TARK, SAT, etc.
- Lucrări in pregătire in curs de trimitere JAIR, Artificial Intelligence Journal, o anumită conferință cu deadline-ul in 17 mai
- Sunt interesat de colaborare pe teme de teoria algoritmicã a jocurilor, agenţi, metode (neuro)simbolice, etc.

# Ce caută un teoretician la un Workshop de Inteligență Artificială?

#### (Submitted on 21 jun 2823)

Large Language Models Still Can't Plan (A Benchmark for LLMs on Planning and Reasoning about Change)

Karthik Valmeekam, Alberto Olmo, Sarath Sreedharan, Subbarao Kambhampati

The seast advances in large languages models (LLbb) have brackers due had de navual larguage procession (MLh. Proc. 15 In EAL). He made of the sea preference or stable language stack is large public forward with early reference impage model. Along with neural language dablies, telve has been a significant interest in unfocutating whether such models, insend on attentions amounted staff, excited measuring callestimes, here the sea has been interest of between data processing bettermore, for various recoming tasks and the preference way in the processing bettermore, for various recoming tasks and the preference way in the preference way the section staff, see an invest bettermore, and extracted processing in the preference way the section staff, seeling the section of the preference way the section staff, seeling the section of the preference way the section staff, seeling the section of the preference way the section staff, seeling the section of the preference way the section staff, seeling the section of the section sectio

Solyson: Computation and Language ResCL); Artificial Intelligence (cs.Al) Clin acc. artist 23 (0.1040); SecCL) (or 200+2200-104004) SecCL) for this ventions

- Agenţii inteligenţi trebuie să interacţioneze cu oamenii şi să inţeleagă raţionalitatea lor, atâta câtă e (sau lipsa ei)
- Proiect (pe termen lung): să incorporăm tot ce ne spun Economia Comportamentală/teoria deciziei despre comportamentul uman intr-o bibliotecă software.

# **Game Theory**



- My payoff may depend not only on what I am doing but also on what others are doing.
- Search for "equilibrium" points.

In this talk: shorthand for Decision Theory + Game Theory.

# **Game Theory in one minute**

$$\begin{array}{c|cc}
C & D \\
C & 2,2 & 0,3 \\
D & 3,0 & 1,1
\end{array}$$

- agents: want to maximize their utility
- Nash equilibrium: action profile so that everyone best-responds to what others are doing.
- may not exist, mixed strategies guaranteed to exist (Nash) by non-constructive methods (Borsuk-Ulam or similar fixed-point theorems).
- Prisoners' Dilemma Paradox: agents will find rational to defect even though if they cooperated they would both be better off

# What's wrong with Game Theory?

#### In Theory:

- Omniscient.
- Perfectly rational
- Sociopath: only thing that matters - its utility.

### In Reality:

- Bounded (if at all) rational.
- Dual process reasoning: fast/heuristic System I, slow/expensive Syst. II.
- Emotional.
- Social.

Richard Thaler called these extremes "Econs" and "Humans".

Normative rather than descriptive: How the Econ <u>should</u> play, rather than how the Human actually plays.

## Why is the Classical Game-Theoretic Person a Caricature?

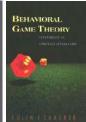
- We're (too much) in love with proving theorems.
- Mathematically tractable models of behavior: simplified.
   Real life: complex/messy.
- Insistence on mathematical (rather than computational) models: we may not even have the "right" framework!

Disclaimer: I will show you some theorems later. In my defense, they all try to point towards what (and how) to implement (or not) game theoretic concepts.

# (Some) Good News



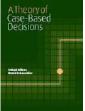


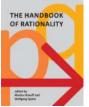










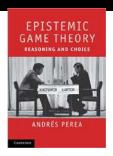


#### **Bad News**

- All this knowledge: far from being implemented/available for building intelligent agents.
- Game theory software: **GAMBIT, nashpy** track classical game theory.
  - Not even clear we have the "right" models to implement!

# OK, so let's use logic to formalize games ...





- ... Except that logic may have many of the same problems.
- Barto, Smets & Solaki (Ertkenntnis, 2021): "Now just as mainstream economics has forgotten Humans to focus on Econs, so has mainstream logic forgotten them to focus on Logons. We name this way the ideal agents studied in "static" epistemic logic with possible worlds semantics (Hintikka 1962) and in AGM belief revision theory (Alchourron et al. 1985). These agents are logically omniscient: perfectly consistent, closed under classical logical consequence in their beliefs, and free from framing effects in their belief revision policies [...]. In fact, Econs may just be Logons engaged in rational choice. The focus on Logons has opened a rift between logic and cognition, similar to the one between the latter and economics"

# Good News: The Cognitive Turn in Logic & KR

Mercier & Sperber: reasoning **not** primarily useful for infer new facts, but to **win arguments** (convince opponents).







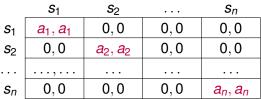
#### Also:

- "natural logic" (Moss): small, natural, tractable fragments.
- "cognitive logics" (Kern-Isberner, see e.g. KR'2021 tutorial)
- cost of reasoning (Solaki & Smets, WOLLIC 2015)
- dual-process logics (Barto, Solaki & Smets, Erkenntnis'19).
  - Bad News: Not that many implementations.
     (Perhaps) reason DL won against Symbolic AI: lots of computational tools, easy to tinker with

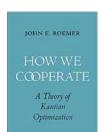
## Now for sampler of concrete (theoretical) results (TARK 2021)

### Kantian equilibria:

 Non-Nashian equilibrium notion, clear(est) definition: symmetric coordination games.



- (loosely) based on Kantian categorical imperative: "act only according to that maxim whereby you can, at the same time, will that it should become a universal law"
- all agents play **the same action**  $x_{OPT}$ , chosen to maximize  $\pi(x, x, ..., x)$ .



# Example: Kantian Equilibria in Prisoners' Dilemma

- Nash equilibrium: both agents defect.
- Kantian equilibrium: both cooperate and are both better off

### Requirements

Interested in minimal rationality equilibrium notions, easily implemented in agents that reason "like people".

#### Therefore, want notions that are:

- expressive: can model plausible scenarios.
- computationally tractable: computing equilibria should be easy (cf van Rooij, The tractable cognition thesis, Cognitive science, 2008).
- cognitively tractable, easy to formally specify: no costly epistemic assumptions (common knowledge, many rounds of iterated elimination of dominated strategies; see also complexity notions of Solaki & Smets)

Main research problem: Are Kantian equilibria (suitably generalized) such a notion? Are there more suitable related/competing notions?

### **Questions & Intuitive Conclusions**

- What is the complexity of (mixed) Kantian equilibria?

NP-hard even for two-player symmetric games. Such equilibria problematic if multiple equilibria exist.

- Can one define Kantian equilibria in more general games?

Yes, in games with certain "symmetry".

- Can one "interpolate" between Kantian and Nash behavior? yes.
- Kantian equilibria or other related notions?

Theoretically interesting, but probably <u>team reasoning</u> more useful for implementations.

# Mixed Kantian equilibria: example

probabilistic combinations of pure strategies. Want combination that maximizes expected payoff when played by everyone

"Platonia Dilemma" (Hofstadter): each agent one of two strategies *S*, *N*. Payoff: 1 if the only agent to play *S*, o otherwise.

- Pure Kantian equilibrium: everyone plays *N*, payoff o, or everyone plays *S*, payoff o.
- mixed Kantian equilibrium: everyone independently plays S w.p.  $\frac{1}{n}$ , o otherwise.

# Symmetric Coordination Games: Mixed Kantian equilibria are easy (but useless)

<u>THEOREM</u>: In symmetric coordination games (more generally in diagonally dominant games:) all mixed Kantian equilibria are pure (hence easy to compute)

Proof idea:  $E[p] = \sum_{i,j} a_{i,j} p_i p_j \le max(a_{i,j}) (\sum p_k)^2 = max(a_{k,k}).$ 

Any equilibrium makes this equality. If there were two different actions i, j in the support of p then we would contradict the diagonally dominant hypothesis.

# Mixed Kantian equilibria: "hard" beyond symmetric coordination games

<u>THEOREM</u>: The following problem, **MIXED KANTIAN EQUILIBRIUM**, is NP-hard:

**INPUT:** Two-player symmetric game G, and an aspiration level  $r \in \mathbb{Q}$ .

**TO DECIDE:** Is there a mixed strategy profile  $x = (x_1, ..., x_m)$  such that the utility of every player under common mixed action  $x_1 a_1 + x_2 a_2 + ... + x_m a_m$  is  $\geq r$ ?

Proof Idea: Follows implicitly from results in the literature (for problem QUADRATIC OPTIMIZATION, which turns out to be equivalent): Motzkin & Strauss (1965), computing  $X_{OPT}$  in o/1 two-player symmetric games equivalent to computing MAX-CLIQUE in equivalent graph.

### NP-hard equilibria are cognitively implausible

Cognitive Science 32 (2008) 939–984
Copyright © 2008 Cognitive Science Society, Inc. All rights reserved.
ISSN: 0364-0213 print/1551-6709 online
DOI: 10.1080/03640210801897856

#### The Tractable Cognition Thesis

#### Iris van Rooij

Nijmegen Institute for Cognition and Information, Radboud University Nijmegen

Received 7 May 2007; received in revised form 26 November 2007; accepted 12 December 2007

#### Abstract

The recognition that human minds/brains are finite systems with limited resources for computation has led some researchers to advance the Tractable Cognition thesis: Human cognitive capacities are constrained by computational tractability. This thesis, if true, serves cognitive psychology by constraining the space of computational-level theories of cognition. To utilize this constraint, a precise and workable definition of "computational tractability" is needed. Following computer science tradition, many cognitive scientists and psychologists define computational tractability as polynomial-time computability, leading to the P-Cognition thesis. This article explains how and why the P-Cognition thesis may be overly restrictive, risking the exclusion of veridical computational-level theories from scientific investigation. An argument is made to replace the P-Cognition thesis by the FPT-Cognition thesis an alternative formalization of the Tractable Cognition thesis (here, FPT stands for fixed-parameter

# **Problems with mixed Kantian Equilibria**

	С	D	E
C	5,5	3,6	1,2
D	6,3	4,4	6,3
Ε	2, 1	3,6	5,5

	С	S
С	10,10	100, 200
S	200,100	6,6

- First game: two Kantian equilibria, (C, C), (E, E) playing mixture of them bad.
- Theorem: "price of miscoordination" for Kantian equilibria.
- Second game: as defined, Kantian equilibria "bad".
   Players would like to anticoordinate.

# Program Equilibria

- Tennenholtz (*Games Econ. Behavior, 2004*): to any game one associate extended game, whose actions are programs.
- agents know the text of other's programs, can act on it.
- program equilibrium = Nash equilibrium of extended game.

IF (your-program == my-program) THEN Cooperate ELSE Defect

# Kantian Program Equilibria

- want: same idea (Kantian equilibrium of extended game).
- only possible when players have identical action sets.
- also: Not clear what a program being "best for all" means in general.
- Tennenholtz's formalization of programs ⇒ paradoxical results.

# Kantian Program Equilibria: Platonia Dilemma

- <u>Platonia Dilemma:</u> protocol "best for all": (collectively) choose a random participant, it plays S, others N.
- can be implemented by agents playing the same program:

Choose a random  $x_i \in \mathbb{Z}_n$ . Broadcast it. Collectively compute  $x = \bigoplus_{k=1}^n x_k$ . If x == i send S, otherwise send N.

# Towards Kantian Program Equilibria

	С	S
С	10, 10	100, 200
S	200, 100	6,6

#### Anticoord(i::ID)

Randomly choose bit  $myb \in \{0,1\}$  send myb to the other player as its otherb.

if  $[myb \oplus otherb \equiv i \pmod{2}]$ then play C

else play S

# Kantian (program) equilibria: our notion

- Only for games with a certain symmetry, quantified by group actions.
- expected player utility across an orbit of the action: the same.

#### **Definition**

A game  $\Gamma$  is called *Pareto symmetric* if there exists a group H acting on the **set of Pareto-optimal action profiles** such that

- For every Pareto optimal profile  $a=(a_1,a_2,\ldots,a_n)$  and  $u\in H$  there exists a permutation  $\sigma\in S_n$  such that  $u\cdot a=(a_{\sigma(1)},\ldots,a_{\sigma(n)})$ .
- For every two players  $i \neq j$  and value  $\lambda$

$$|\{u \in H : (u \cdot a)_i = \lambda\}| = |\{u \in H : (u \cdot a)_j = \lambda\}|$$

- notion of program: technical.

# Kantian equilibria in Pareto symmetric games

	C	S
С	10, 10	100,200
S	200, 100	6,6

	C	D
C	2,2	0,3
D	3,0	1,1

Platonia Dilemma: Kantian equilibrium orbit (S, N, ..., N), (N, S, ..., N), ... (N, N, ... S).

 Theorem (extended version): Kantian equilibrium can be characterized as convex combinations of orbit(s) that maximize player expected payoffs.

### **Conclusions/Work in Progress**

- Mixed Kantian equilibria: not cognitively plausible.
- generalization: goes through program equilibria.
  - How do agents recognize symmetry in (our notion of) Kantian games?
- Hopeless if game represented by game matrix
- (In progress) Connect it to general game playing. Version of Game Description Language (GDL, Genesereth et al., AI Magazine 2005, Tielscher IJCAI 2019) with symmetries embedded in the description
  - (in progress) Implementing Kantian optimization in (our version of) GDL.